OS Concepts and structure

Today

- OS services
- OS interface to programmers/users
- OS components & interconnects
- Structuring Oss

Next time

- Processes

OS Services

- OS provides a number of services
  - To users directly through a command interpreter/shell or GUI
  - To application programs through system calls
- Some services are for convenience
  - Program execution
  - I/O operation
  - File system management
  - Communication
- Some to ensure efficient operation
  - Resource allocation
  - Accounting
  - Protection and security

Command interpreter (shell) & GUI

- Command interpreter
  - Handle (interpret and execute) user commands
  - Could be part of the OS
    - MS DOS, Apple II
  - Could be just a special program
    - UNIX, Windows XP
    - In this way, multiple options – shells – are possible
  - The command interpreter could
    - Implement all commands
    - Simply understand what program to invoke and how (UNIX)
- GUI
  - Friendlier, through a desktop metaphor, if sometimes limiting
  - Xerox PARK Alto >> Apple >> Windows >> Linux

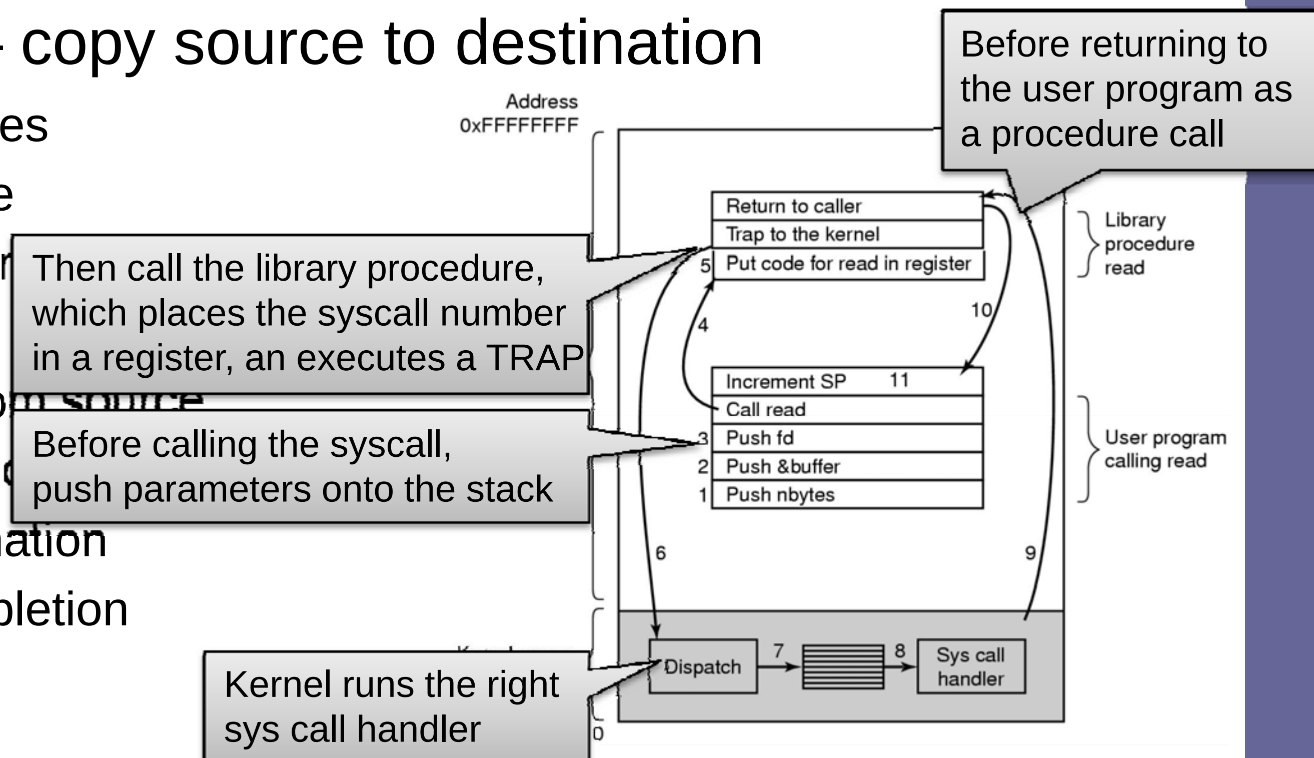
Shell – stripped down

```
while (TRUE) {                          /* repeat forever */
    type_prompt( );                     /* display prompt */
    read_command(command, parameters) /* input from terminal */

if (fork() != 0) {                     /* fork off child process */
    /* Parent code */
    waitpid( -1, &status, 0); /* wait for child to exit */
} else {
    /* Child code */
    execve (command, parameters, 0); /* execute command */
 }
}
```

System calls

- Low-level interface to services for application programs
- Higher-level requests get translated into sequence of system calls
- Writing cp – copy source to destination
  - Get file names
  - Open source
  - Create destir
  - Loop
    - Read fro… source
    - Copy to …
  - Close destination
  - Report completion
  - Terminate



Making the system call: `read(fd, buffer, nbytes)`

Processes

- A program in execution
  - Address space
  - Set of registers
- To get a better sense of it
  - What data do you need to (re-) start a suspended process?
  - Where do you keep this data?
  - What is the process abstraction I/F offered by the OS
    - Create, delete, suspend, resume & clone a process
    - Inter-process communication & synchronization
    - Create/delete a child process

Call	**Description**
pid = fork()	Create a child process identical to the parent
pid = waitpid(pid, &statloc, options)	Wait for a child to terminate
s = execve(name, argv, environp)	Replace a process' core image
exit(status)	Terminate process execution & return status

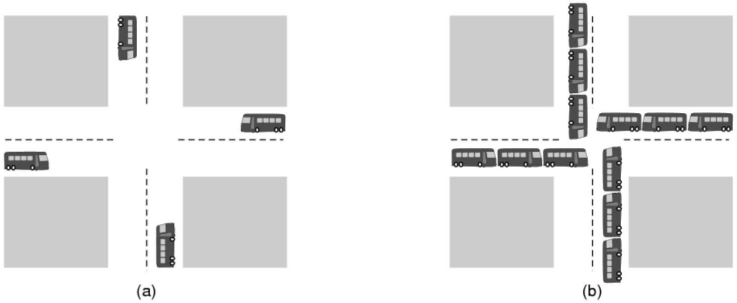
Memory management

- Main memory – the directly accessed storage for CPU
  - Programs must be stored in memory to execute
  - Memory access is fast (e.g., 60 ns to load/store)
    - but memory doesn't survive power failures
- OS must:
  - Allocate memory space for programs (explicitly and implicitly)
  - Deallocate space when needed by rest of system
  - Maintain mappings from physical to virtual memory
    - e.g. through page tables
  - Decide how much memory to allocate to each process
  - Decide when to remove a process from memory

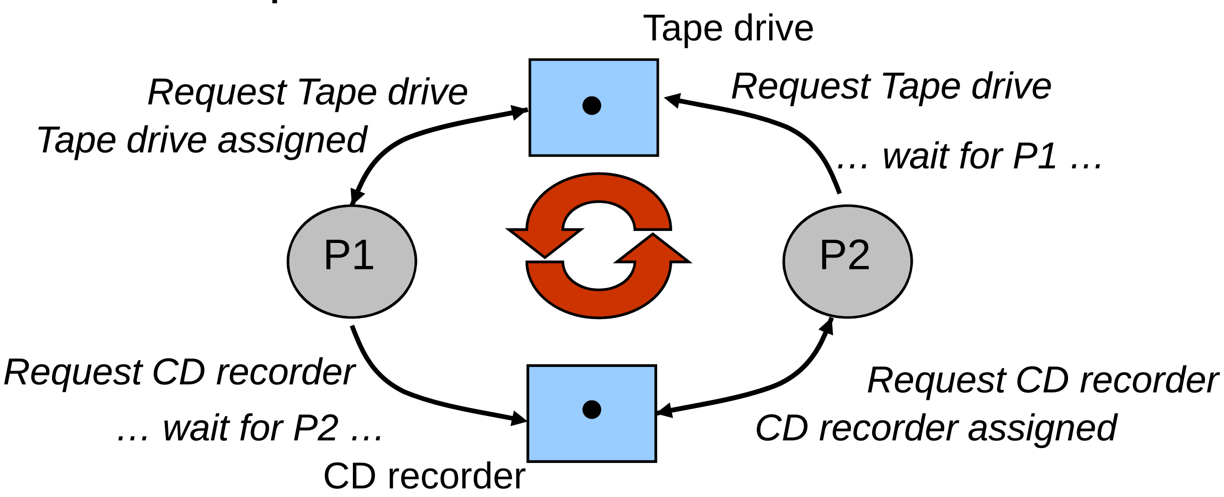
Call	Description
void *sbrk(intptr_t increment)	Increments program data space by 'increment' bytes

Deadlocks

- Processes interactions & deadlock opportunities
- A real-world example



- A simple OS example
  - Two processes making a CD from tape



I/O

- A big chunk of the OS kernel deals with I/O
  - Hundreds of thousands of lines in NT
- The OS provides a standard interface between programs (user or system) and devices
  - file system (disk), sockets (network), frame buffer (video)
- Device drivers are the routines that interact with specific device types
  - Encapsulates device-specific knowledge
    - e.g., how to initialize a device, how to request I/O, how to handle interrupts or errors
  - Examples: SCSI device drivers, Ethernet card drivers, video card drivers, sound card drivers, …

Secondary storage

- Secondary storage (disk, tape) is persistent memory
  - Often magnetic media, survives power failures (hopefully)
- Routines that interact with disks are typically at a very low level in the OS
  - Used by many components (file system, VM, …)
  - Handle scheduling of disk operations, head movement, error handling, and often management of space on disks
- Usually independent of file system
  - Although there may be cooperation
  - File system knowledge of device details can help optimize performance
    - e.g., place related files close together on disk

File systems

- Secondary storage devices are hard to work with
- File system offers a convenient abstraction
  - Defines logical abstractions/objects like files & directories
  - As well as operations on these objects
- A file is the basic unit of long-term storage
  - File: named collection of persistent information
- A directory is just a special kind of file
  - Directory: file containing names of other files & metadata
- Interface:
  - File/directory creation/deletion, manipulation, copy, lock
- Other higher level services: accounting & quotas, backup, indexing or search, versioning

System calls

File management

Call	Description
fd = open(file, how, …)	Open a file for reading, writing or both.
s = close(fd)	Close an open file
n = read(fd, buffer, nbytes)	Read data from a file into a buffer
n = write(fd, buffer, nbytes)	Write data from a buffer into a file
pos = lseek(fd, offest, whence)	Move the file pointer
s = stat(name,&buf)	Get a file’s status info

Directory & file system management

Call	Description
s = mkdir(name, mode)	Create a new directory
s = rmdir(name)	Remove an empty directory
s = link(name1, name2)	Create a new entry, name2, pointing to name1
s = unlink(name)	Remove a directory entry
s = mount(special, name, flag)	Mount a file system
s = unmount(special)	Unmount a file system

Protection

- Protection is a general mechanism used throughout the OS
  - All resources needed to be protected
    - memory
    - processes
    - files
    - devices
    - …
  - Protection mechanisms help to detect and contain errors, as well as preventing malicious destruction

OS design & implementation

- A design task – start from goals & specification
- Affected by choice of hardware, type of system
- *User* goals and *System* goals
  - User – convenient to use, easy to learn, reliable, safe, fast
  - System – easy to design, implement, & maintain, also flexible, reliable, error-free & efficient
- Clearly conflicting goals, no unique solution
- Some other issues complicating this
  - Size: Windows XP ~40G SLOC, RH 7.1 17G SLOC
  - Concurrency – multiple users and multiple devices all active at once
  - Potentially hostile users, but some users want to collaborate
  - Long expected lives & no clear ideas on how the system will be used
  - Portability and support to thousands of device drivers
  - Backward compatibility

OS design & implementation

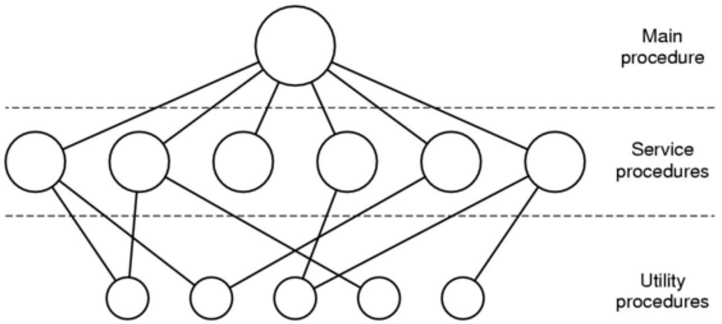
- A software engineering principle – separate policy & mechanism
  - **Policy:** What will be done?
  - **Mechanism:** How to do it?
  - Why do you care? Maximum flexibility, easier to change policies
- Implementation on high-level language
  - Early on – assembly (e.g. MS-DOS – 8088), later Algol (MCP), PL/1 (MULTICS), C (Unix, …)
  - Advantages – faster to write, more compact, easier to maintain & debug, easier to port
  - Cost – Slower, but who cares?!

OS structure

- OS made of number of components
  - Process management, memory management, file system, …
  - and System programs (privileged and non-privileged)
    - e.g., bootstrap code, the init program, …
- Major design issue
  - How do we organize all this?
  - What are all of the code modules, and where do they exist?
  - How do they interact?
- Massive software engineering and design problem
  - design a large, complex program that:
    - performs well, is reliable, is extensible, is backwards compatible, …

Monolithic design

- Major advantage:
  - Cost of module interactions is low (procedure call)
- Disadvantages:
  - Hard to understand
  - Hard to modify
  - Unreliable (no isolation between system modules)
  - Hard to maintain
- What is the alternative?
  - Find a way to organize the OS in order to simplify its design and implementation



Layering

- The traditional approach is layering
  - Implement OS as a set of layers
  - Each layer presents an enhanced ‘virtual machine’ to the layer above
- Each layer can be tested and verified independently (Dijkstra’s THE system)

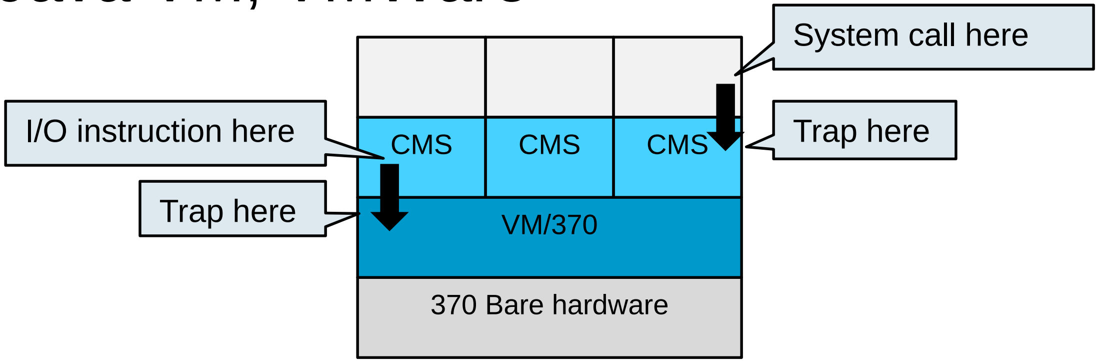
Layer	**Description**
5: Job managers	Execute users’ programs
4: Device managers	Handle device & provide buffering
3: Console manager	Implements virtual consoles
2: Page manager	Implements virtual memory for each process
1: Kernel	Implements a virtual processor for each process
0: Hardware	

Problems with layering

- Imposes hierarchical structure
  – but real systems are more complex:
    • File system requires VM services (buffers)
    • VM would like to use files for its backing store
  – Strict layering isn't flexible enough
- Poor performance
  – Each layer crossing implies overhead
- Disjunction between model and reality
  – Systems modeled as layers, but not really built that way

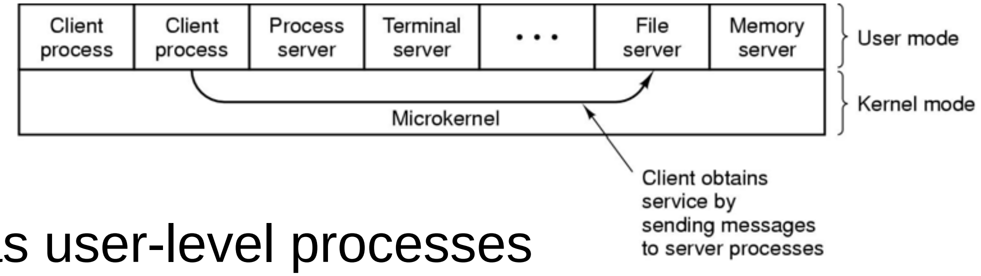
Virtual machines

- Initial release of OS/360 were strictly batch but users wanted timesharing
  – IBM CP/CMS, later renamed VM/370 (‘79)
- Observation – Timesharing system provides (1) multiprogramming & (2) extended (virtual) machine
- Essence of VM/370 – separate the two
  – Heart of the system (VMM) does multiprogramming & provides to next layer up multiple exact copies of bare HW
  – Each VM can run any OS
- More recently – Java VM, VMWare



Microkernels

- Popular in the late 80's, early 90's
  - Recent resurgence
- Goal:
  - Minimize what goes in kernel
  - Organize rest of OS as user-level processes



- This results in:
  - Better reliability (isolation between components)
  - Ease of extension and customization
  - Poor performance (user/kernel boundary crossings)
- First microkernel system was Hydra (CMU, 1970)
  - Follow-ons: Mach (CMU), Chorus (French UNIX-like OS), OS X (Apple), in some ways NT (Microsoft)

Operating system generation

- OS design for a class of machines; need to configure it for yours - SYSGEN
- SYSGEN program gets info on specific configuration
  - CPU, memory, devices, other parameters
- Once you got it you could
  - Modify source code & recompile kernel
  - Modify tables and select precompiled modules
  - Modify tables; everything is there & selection is at run time

  Trading size & generality for ease of modification

System boot

How does the OS gets started?

- Booting: starting a computer by loading the kernel
- Instruction register loaded with predefined memory location – bootstrap loader (ROM)
  - Why not just put the OS in ROM? Cell phones & PDAs
- Bootstrap loader
  - Run diagnostics
  - Initialize registers & controllers
  - Fetch second bootstrap program form disk
    - Why do you need a second bootstrap loader?
- Second bootstrap program loads OS & gets it going
  - A disk with a boot partition – boot/system disk

Summary & preview

- Today
  - The mess under the carpet
  - Basic concepts in OS
  - Structuring OS - a few alternatives

- Next …
  - Process – the central concept in OS
    - Process model and implementation
  - Threads – a light-weight process
    - Thread model, usage & implementation